

BEAR: sketching BFGS algorithm for ultra-high dimensional feature selection in sublinear memory

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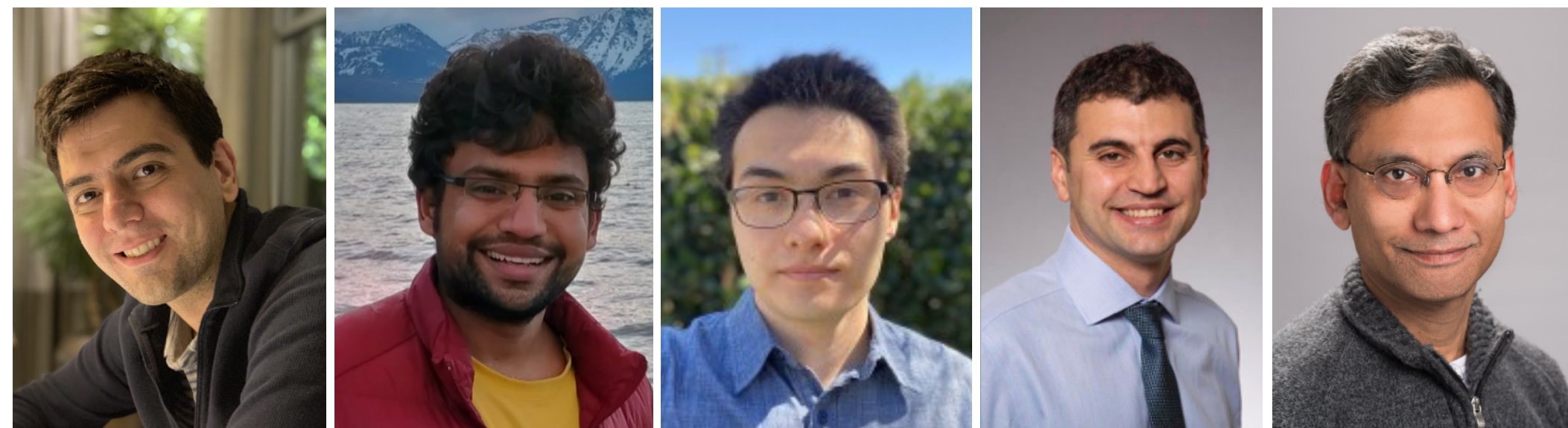
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big and high dimensional data in everyday life

- web services
 - language processing
 - networking
 - genomics/proteomics
 - health-care
-
- critical need for **scalable algorithms** to extract **important features** from the data
 - limited computing **resource**



problem setup

- n data points $(\theta_i)_{i=1}^n = (\mathbf{x}_i, y_i)_{i=1}^n$
living in ultra-high dimensional feature space $\mathbf{x}_i \in \mathbb{R}^p$ ($p \gg n$)
goal: find a **small subset of features** best explains the output

>10¹⁵!

- k -**sparse** feature vector $\beta^* \in \mathbb{R}^p$
loss function $f(\beta, \theta) : \mathbb{R}^p \rightarrow \mathbb{R}$

challenge: not enough **memory** to store the **intermediately dense** feature vector β (sublinear alg.)

- optimization problem $\min_{\beta} \sum_{i=1}^n f(\beta; \theta_i)$

- stochastic gradient descent (SGD) $\beta_{t+1} = \beta_t - \eta_t \mathbf{g}(\beta_t; \Theta_t)$
minibatch $\Theta_t = \{\theta_{t1}, \theta_{t2}, \dots, \theta_{tb}\}$
with the SGD term defined as $\mathbf{g}(\beta_t; \Theta_t) = \sum_{i=1}^b \nabla_{\beta_t} f(\beta_t; \theta_{ti})$

Count Sketch (CS)

- data structure to **compressively** store the **number of occurrences** of many number of streaming items

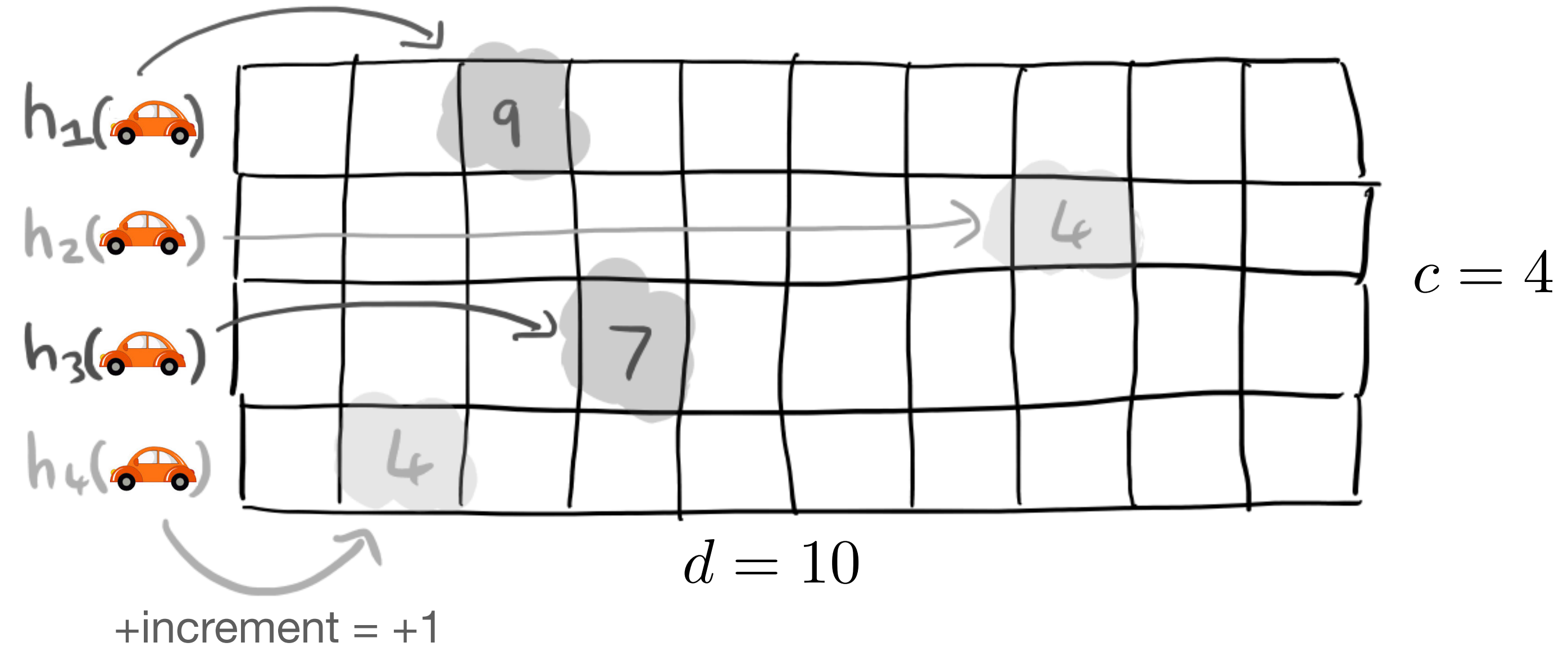
- fast operations

- ADD (item, increment)
- QUERY (item)

🚗 $\approx \text{median}(\{4, 4, 7, 9\})$

random hash function

$$h_j : \{1, 2, \dots, p\} \rightarrow \{1, 2, \dots, d\}$$



items (p)

all colors

$$m = d \times c$$

memory of CS

frequent items (k)

top colors



Count Sketch (CS)

- data structure to **compressively** store the **number of occurrences**

random hash function

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Theorem 1 *Charikar et al. (2002)* Count Sketch finds *top-k items* z_i with $\pm \epsilon \|\mathbf{z}\|_2$ error, with probability at least $1 - \delta$, in space $\mathcal{O}(\log(\frac{p}{\delta})(k + \frac{\|\mathbf{z}^{tail}\|_2^2}{(\epsilon\zeta)^2}))$, where $\|\mathbf{z}^{tail}\|_2^2 = \sum_{i \notin top-k} z_i^2$ is the energy of the non-top-k items and ζ is the k^{th} largest value in \mathbf{z} .

- ADD (item, increment)
- QUERY (item)

🚗 $\approx \text{median}(\{4, 4, 7, 9\})$

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frequent items (k)

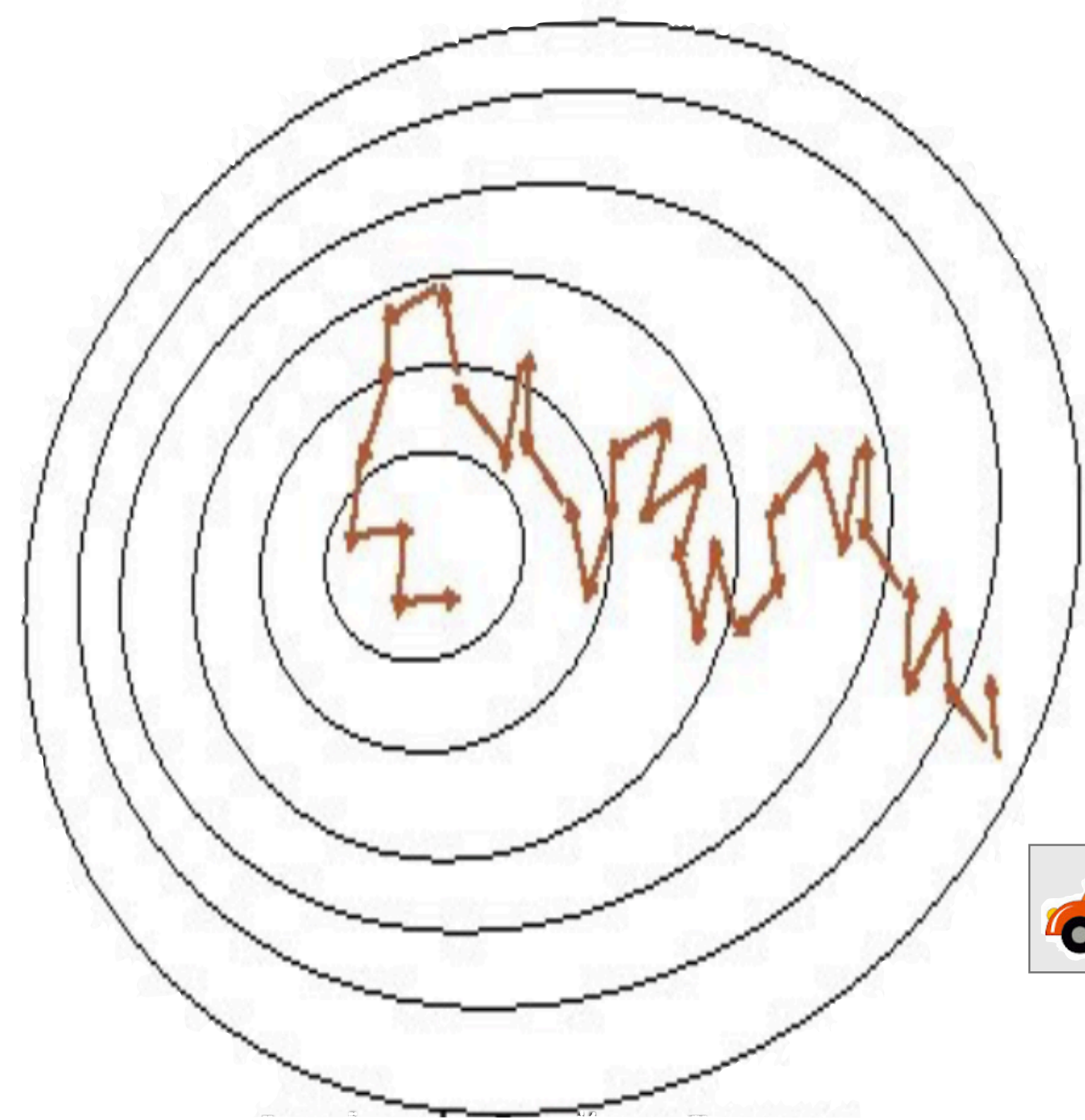
top colors




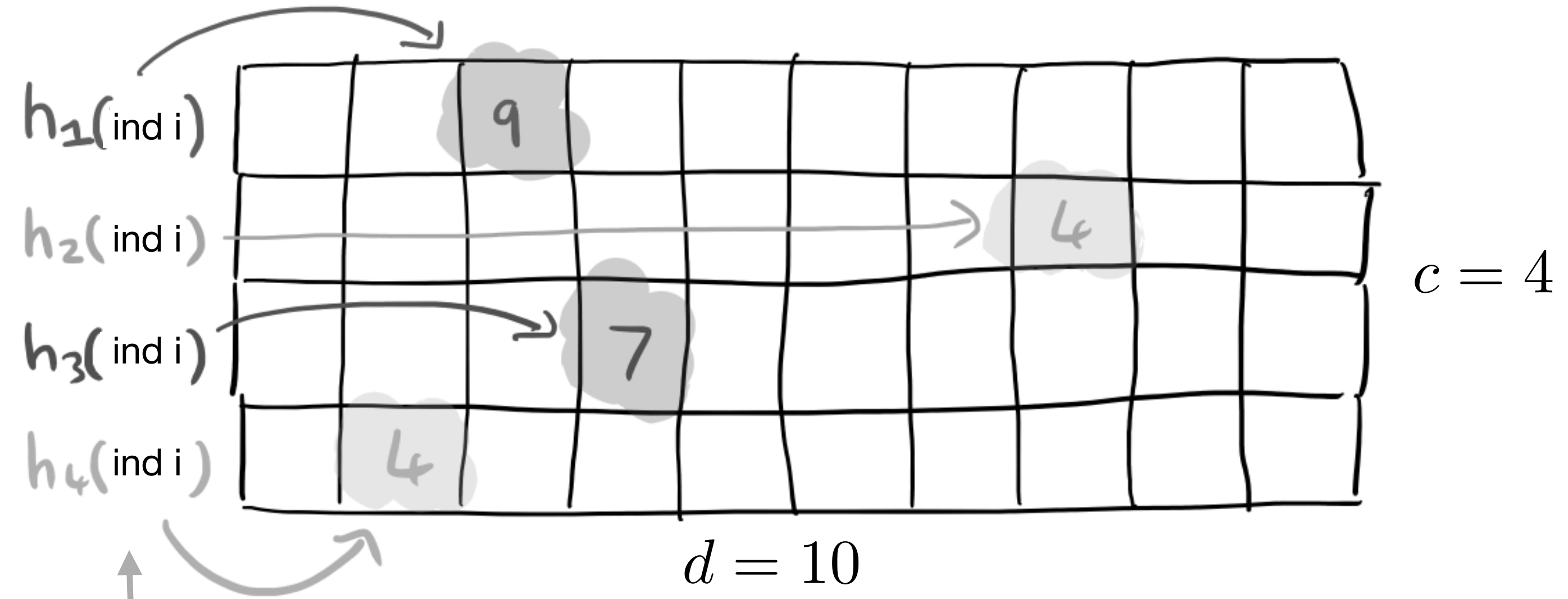
feature selection with CS

$$h_j : \{1, 2, \dots, p\} \rightarrow \{1, 2, \dots, d\}$$

$$\beta_{t+1} = \beta_t - \eta_t \mathbf{g}(\beta_t; \Theta_t)$$



 $\equiv \eta_t \mathbf{g}_i(\cdot)$



+increment = $-\eta_t \mathbf{g}_i(\cdot)$

items (p)

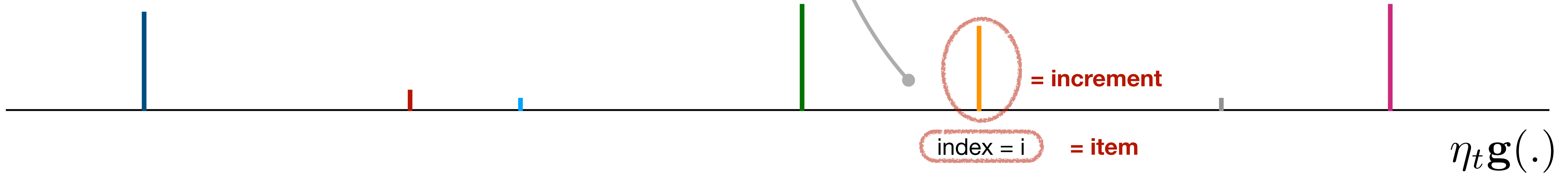
all **features**

$$m = d \times c$$

memory of CS

frequent items (k)

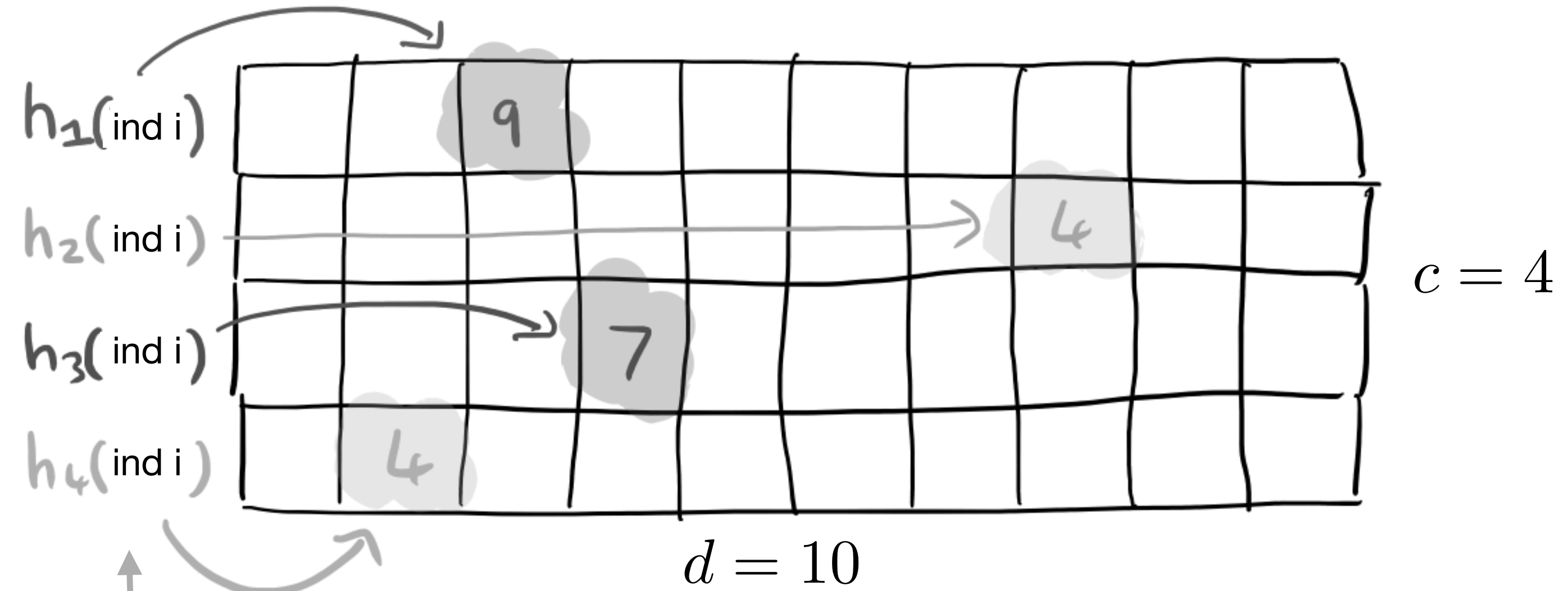
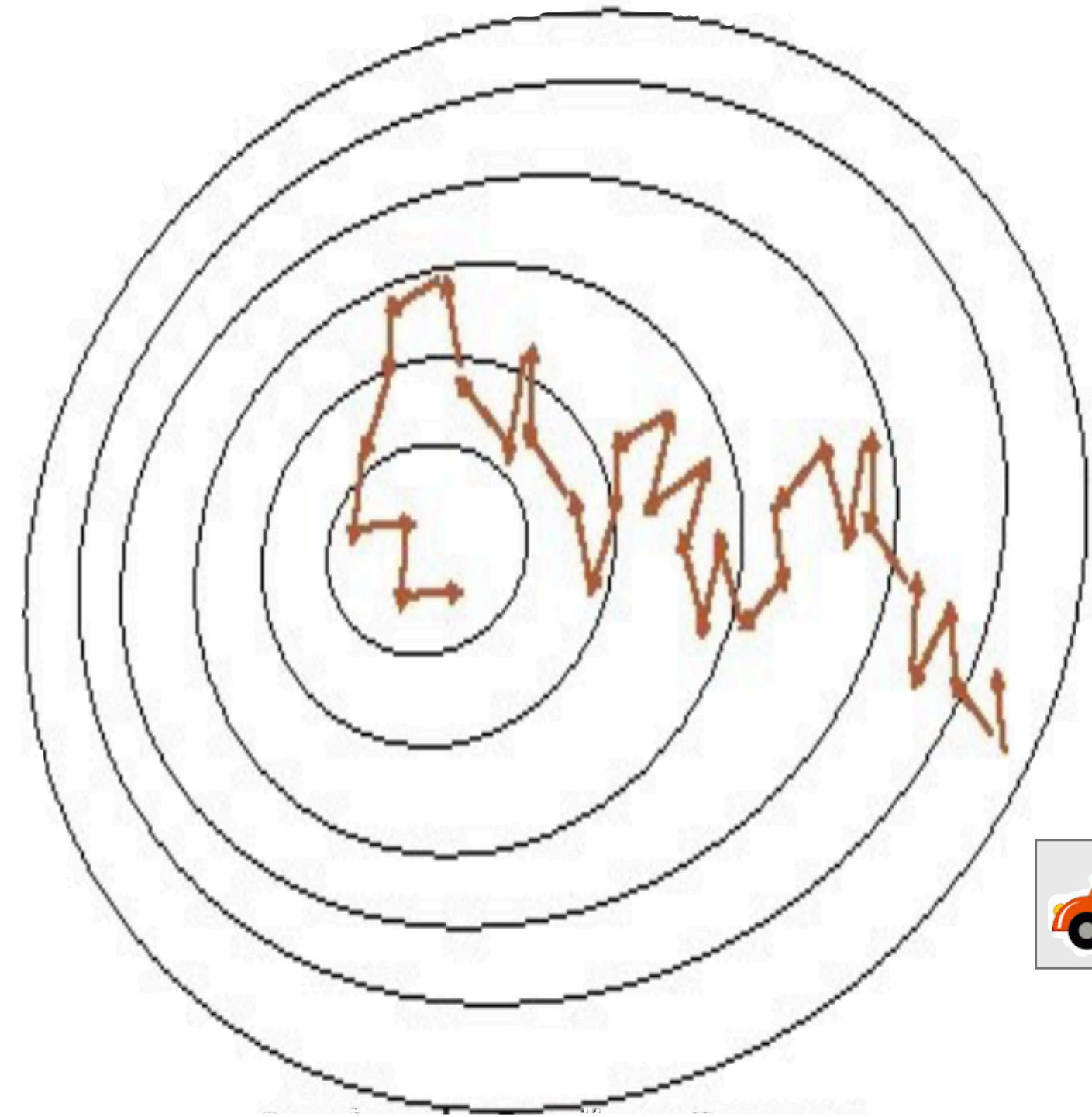
top **features**



feature selection with CS

$$h_j : \{1, 2, \dots, p\} \rightarrow \{1, 2, \dots, d\}$$

MISSION : $\beta_{t+1}^s = \beta_t^s - \eta_t \mathbf{g}^s(\text{Query}_{\text{top-}k}(\beta_t^s); \Theta_t)$



+increment = $-\eta_t \mathbf{g}_i(\cdot)$

items (p)

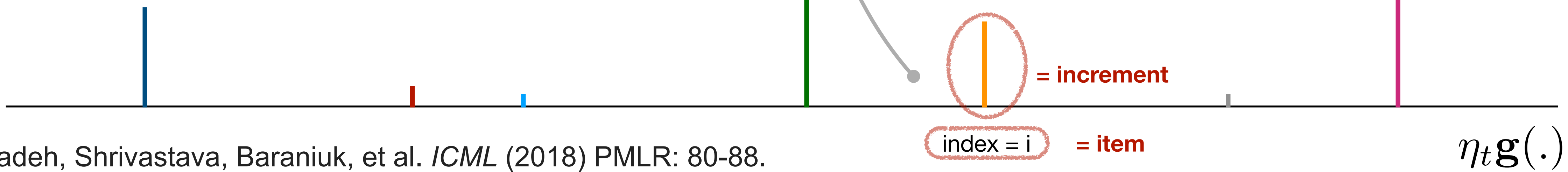
all **features**

$$m = d \times c$$

memory of CS

frequent items (k)

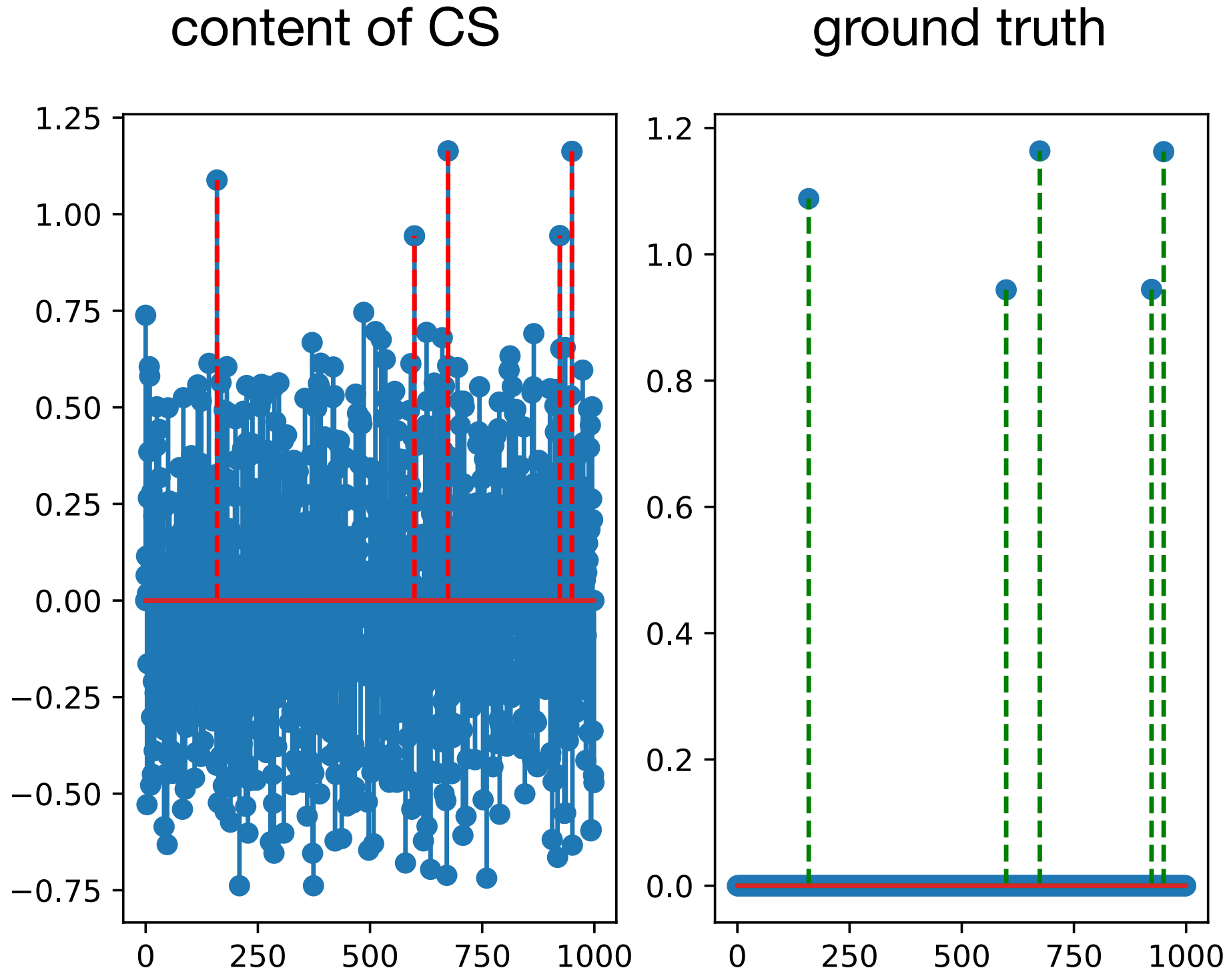
top **features**



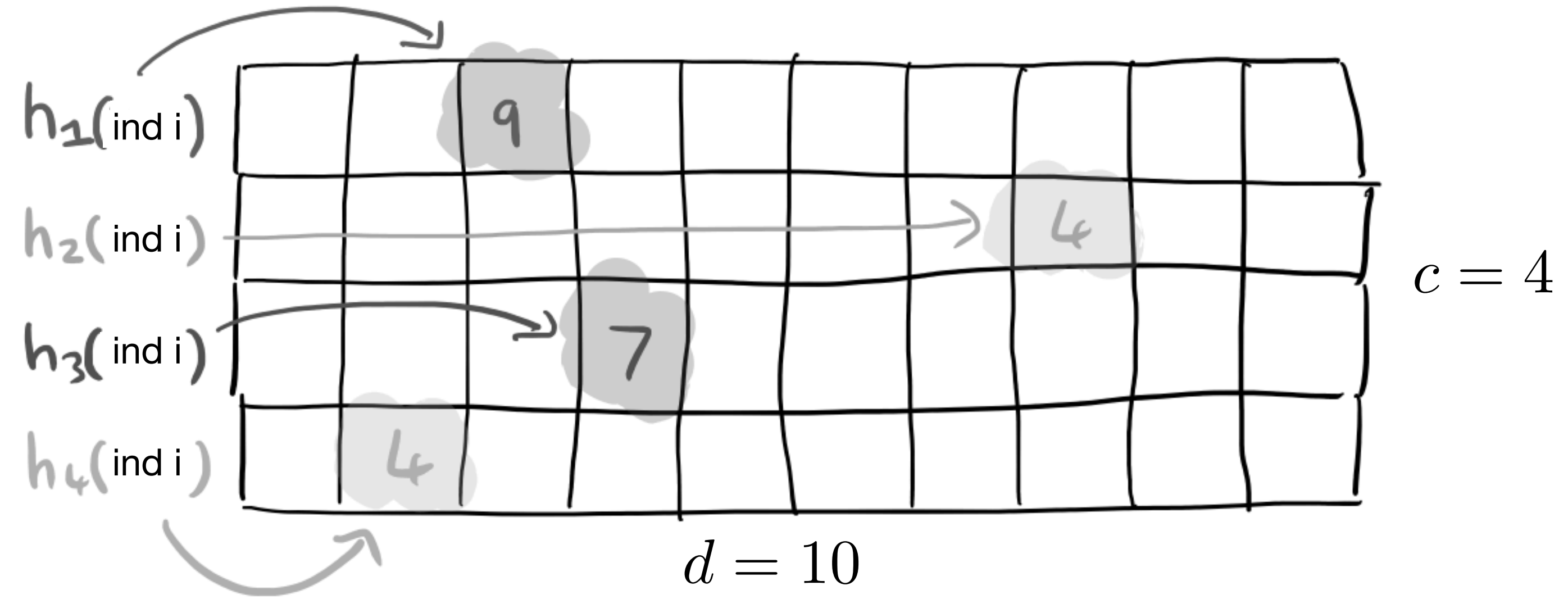
feature selection with CS

MISSION : $\beta_{t+1}^s = \beta_t^s - \eta_t \mathbf{g}^s(\text{Query}_{\text{top-}k}(\beta_t^s); \Theta_t)$

after convergence



$$h_j : \{1, 2, \dots, p\} \rightarrow \{1, 2, \dots, d\}$$



observation: sketch of noisy component of SGD in CS do not cancel out and results in memory wasted to store sketched noise

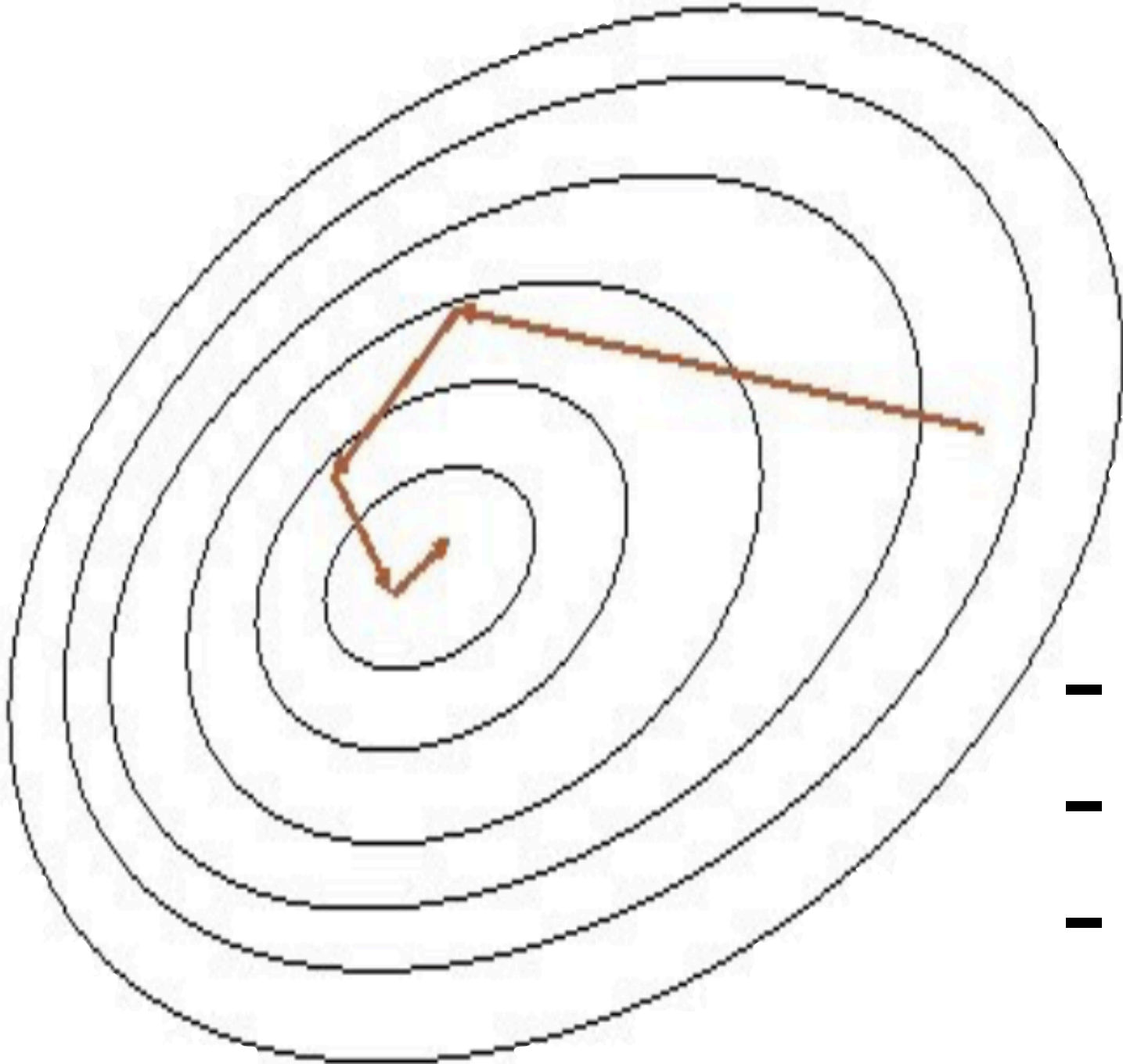
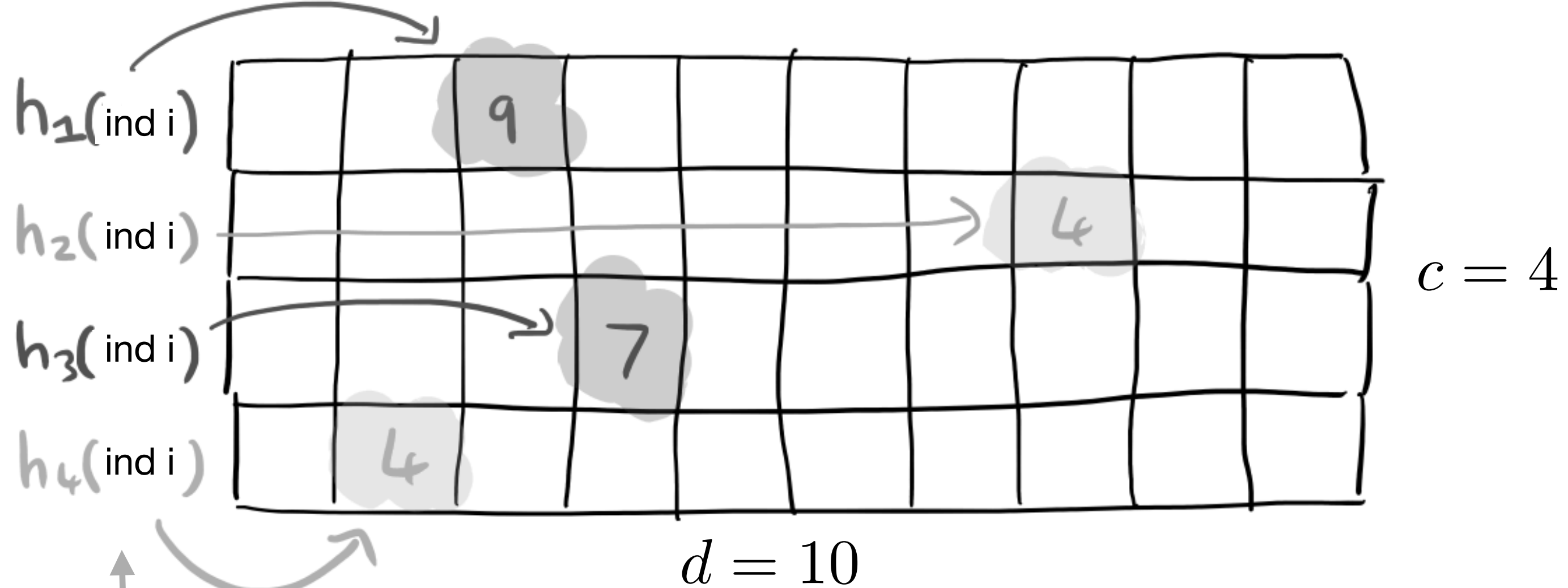
Theorem 1 Charikar et al. (2002) Count Sketch finds top-k items z_i with $\pm \epsilon \|\mathbf{z}\|_2$ error, with probability at least $1 - \delta$, in space $\mathcal{O}(\log(\frac{p}{\delta})(k + \frac{\|\mathbf{z}^{\text{tail}}\|_2^2}{(\epsilon \zeta)^2}))$, where $\|\mathbf{z}^{\text{tail}}\|_2^2 = \sum_{i \notin \text{top-}k} z_i^2$ is the energy of the non-top-k items and ζ is the k^{th} largest value in \mathbf{z} .

idea: second order sketching

$$h_j : \{1, 2, \dots, p\} \rightarrow \{1, 2, \dots, d\}$$

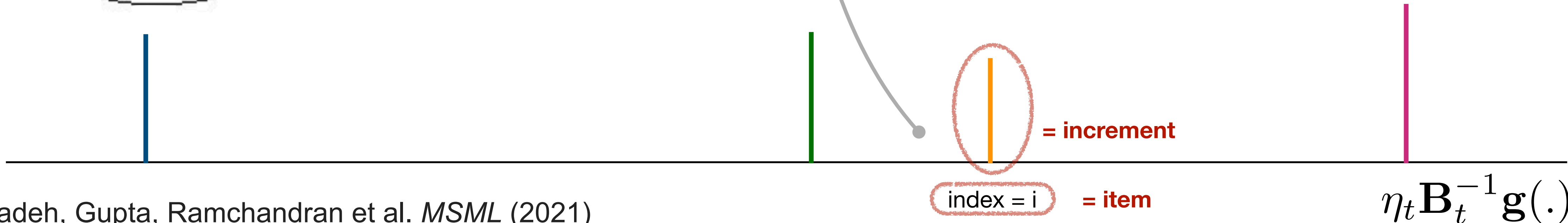
$$\beta_{t+1} = \beta_t - \eta_t \mathbf{B}_t^{-1} \mathbf{g}(\beta_t, \Theta_t)$$

$$\mathbf{B}_t = \nabla_{\beta_t}^2 f(\beta_t, \Theta_t) \in \mathbb{R}^{p \times p}$$



- more comp. cost per iteration
- less noisy gradient
- memory-accuracy tradeoff

question: how to compute/store the Hessian?

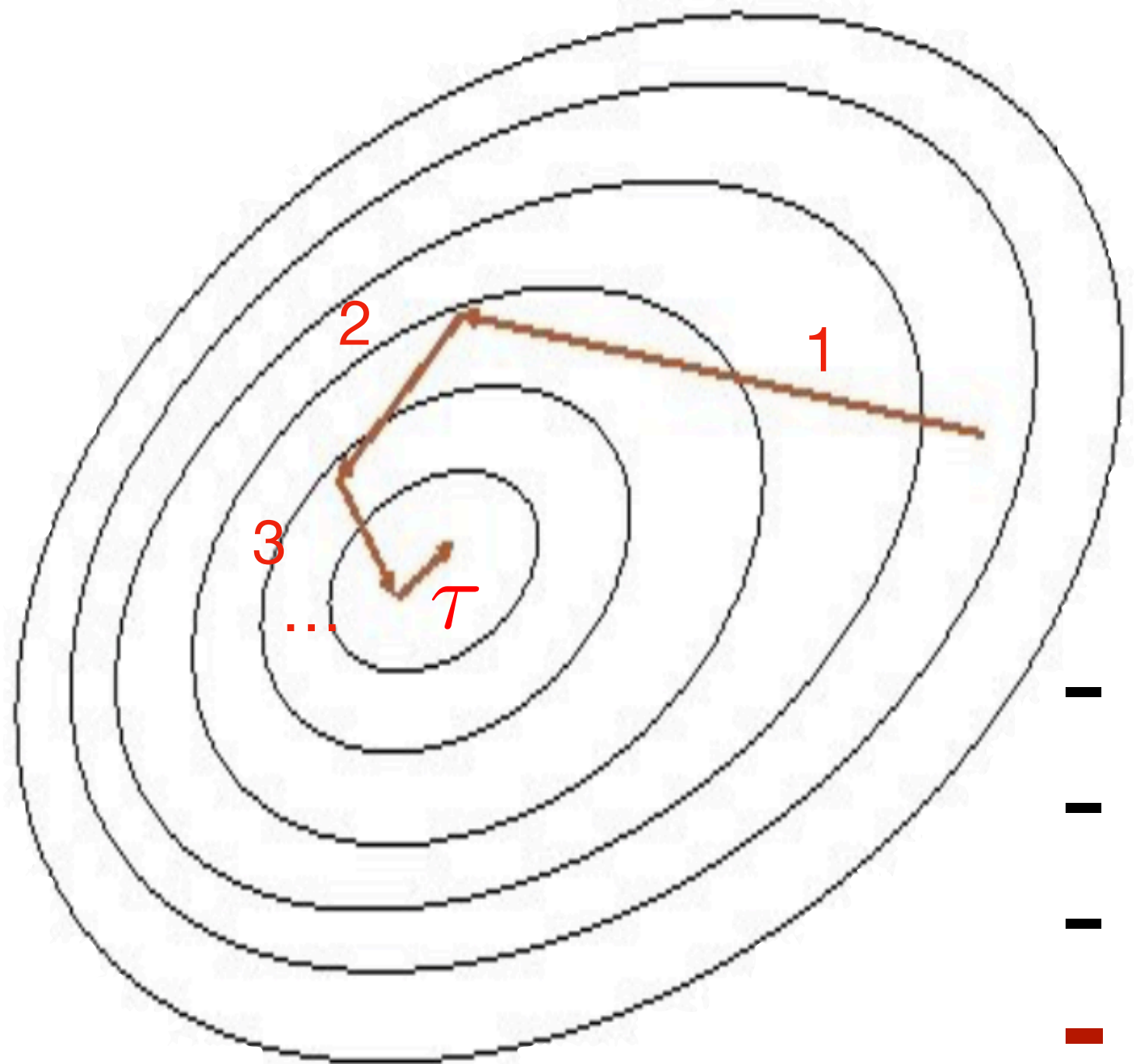


limited-memory BFGS

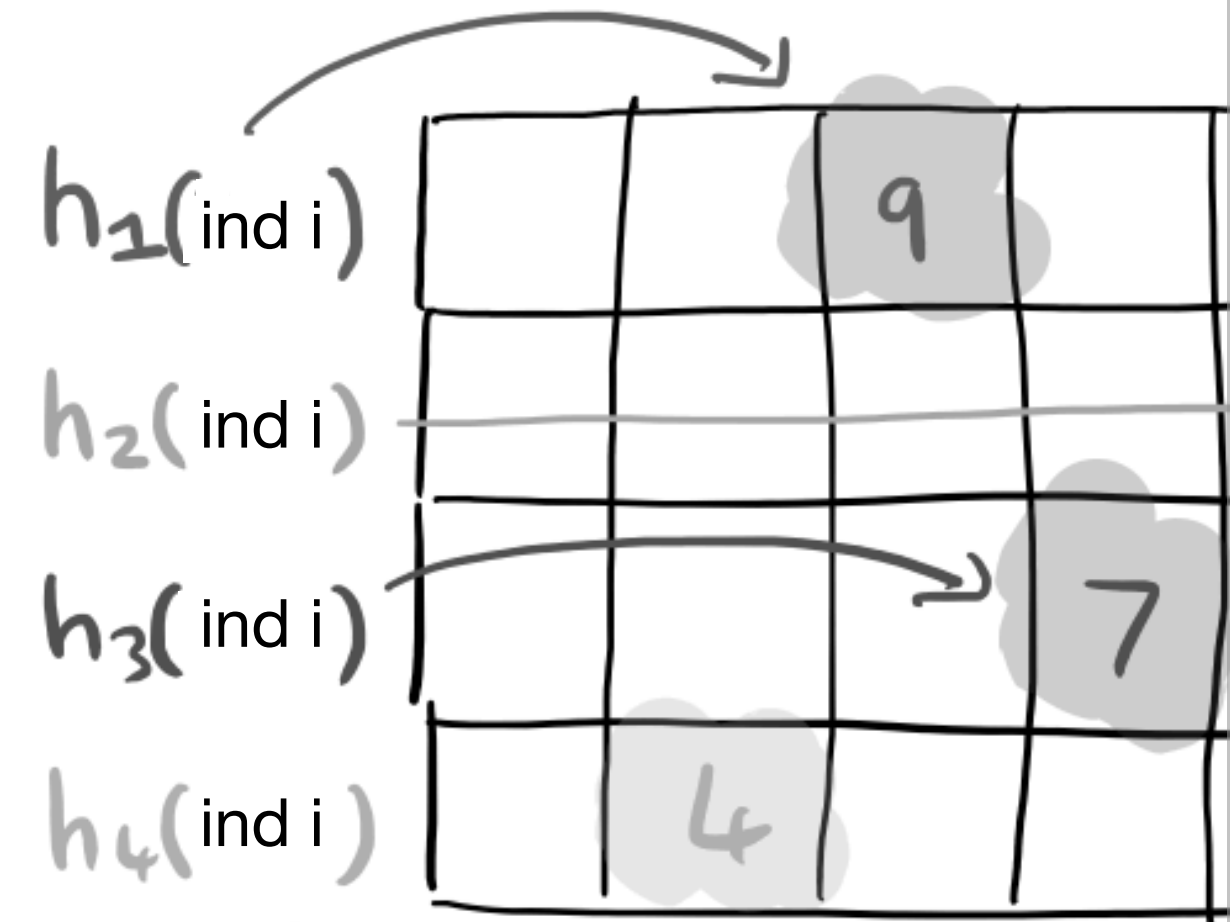
$$h_j : \{1, 2, \dots, p\} \rightarrow \{1, 2, \dots, c\}$$

$$\beta_{t+1} = \beta_t - \eta_t \mathbf{B}_t^{-1} \mathbf{g}(\beta_t, \Theta_t)$$

$$\mathbf{B}_t = \nabla_{\beta_t}^2 f(\beta_t, \Theta_t) \in \mathbb{R}^{p \times p}$$



- more comp. cost per iteration
- less noisy gradient
- memory-accuracy tradeoff
- **no need to store/compute inverse Hessian**



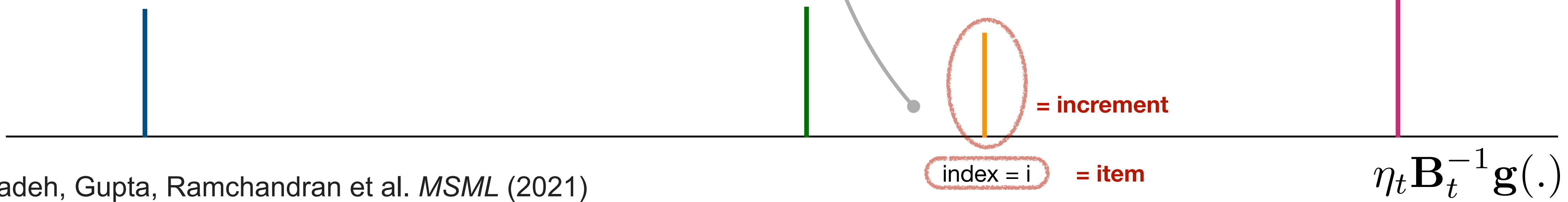
Algorithm 1 Limited-memory BFGS

Input: $\mathbf{g}(\hat{\beta}_t, \Theta_t)$ and $\{\mathbf{s}_i, \mathbf{r}_i\}_{i=t-\tau+1}^t$

1. $\rho_t = \frac{1}{\mathbf{r}_t^T \mathbf{s}_t}$.
2. $\mathbf{q}_t = \mathbf{g}(\hat{\beta}_t, \Theta_t)$,
 for $i = t$ to $t - \tau + 1$:
 $\alpha_i = \rho_i \mathbf{s}_i^T \mathbf{q}_i$,
 $\mathbf{q}_{i-1} = \mathbf{q}_i - \alpha_i \mathbf{r}_i$.
3. $\mathbf{z}_{t-\tau} = \frac{\mathbf{r}_t^T \mathbf{s}_t}{\mathbf{r}_t^T \mathbf{r}_t} \mathbf{q}_{t-\tau}$,
 for $i = t - \tau + 1$ to t :
 $\gamma_i = \rho_i \mathbf{r}_i^T \mathbf{z}_i$.
 $\mathbf{z}_i = \mathbf{z}_{i-1} + \mathbf{s}_i (\alpha_i - \gamma_i)$.

Return: \mathbf{z}_t

**approximate $\mathbf{B}_t^{-1} \mathbf{g}(\cdot)$ using
gradients from last few
 τ iterations**



BEAR algorithm: sketch LFBGS gradients using CS

find the descent
direction using LFBGS
and update CS

Algorithm 2 BEAR

Initialize: $t = 0$, Count Sketch $\beta_{t=0}^s = 0$, top- k heap.
while stopping criteria not satisfied **do**

1. Sample b independent data points in a minibatch $\Theta_t = \{\theta_{t1}, \dots, \theta_{tb}\}$.
2. Find the active set \mathcal{A}_t .

3. QUERY the feature weights in $\mathcal{A}_t \cap \text{top-}k$ from Count Sketch $\beta_t = \text{query}(\beta_t^s)$.
4. Compute stochastic gradient $\mathbf{g}(\beta_t, \Theta_t)$.
5. Compute the descent direction with Alg. 1 $\mathbf{z}_t = \text{LFBGS}(\mathbf{g}(\beta_t, \Theta_t), \{\mathbf{s}_i, \mathbf{r}_i\}_{i=t-\tau+1}^t)$.
6. ADD the sketch of \mathbf{z}_t at the active set $\hat{\mathbf{z}}_t = \mathbf{z}_t^{\mathcal{A}_t}$ to Count Sketch $\beta_{t+1}^s := \beta_t^s - \eta_t \hat{\mathbf{z}}_t^s$.

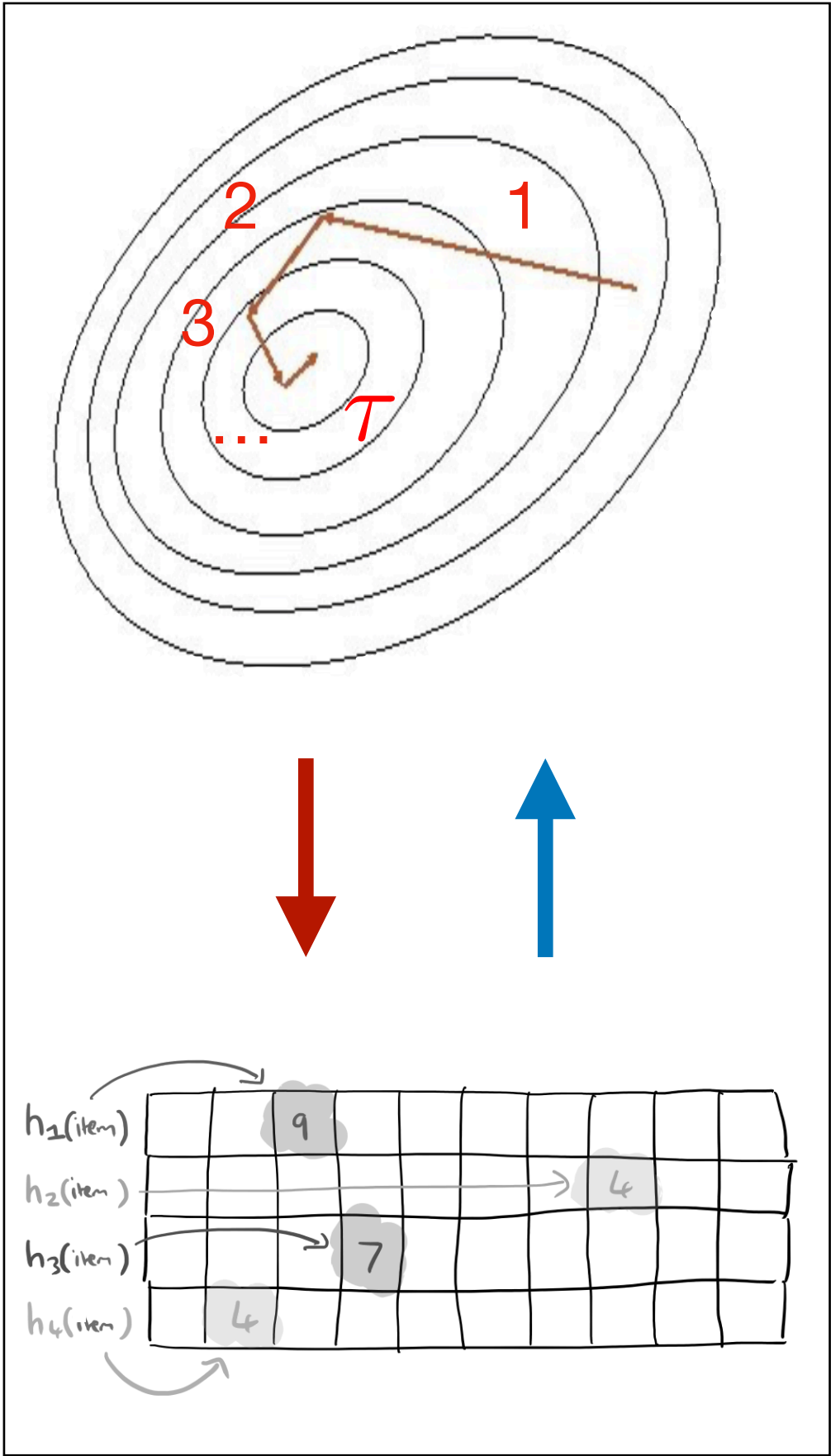
7. QUERY the features weights in $\mathcal{A}_t \cap \text{top-}k$ from Count Sketch $\beta_{t+1} = \text{query}(\beta_{t+1}^s)$.
8. Compute stochastic gradient $\mathbf{g}(\beta_{t+1}, \Theta_t)$.
9. Set $\mathbf{s}_{t+1} = \beta_{t+1} - \beta_t$, and $\mathbf{r}_{t+1} = \mathbf{g}(\beta_{t+1}, \Theta_t) - \mathbf{g}(\beta_t, \Theta_t)$.

10. Update the top- k heap.
11. $t = t + 1$.

end while

Return: The top- k heavy-hitters in Count Sketch.

query CS and store the
gradient and feature
difference vectors



convergence

Theorem 2 *Let $f(\cdot)$ and the step sizes η_t satisfy the assumptions above. Let the size of Count Sketch be $m = \theta(\varepsilon^{-2} \log 1/\delta)$ with number of hashes $d = \theta(\varepsilon^{-1} \log 1/\delta)$ for $\varepsilon, \delta > 0$. Then, the Euclidean distance between updates β_t^s in the BEAR algorithm and the sketch of the solution of problem (1) converges to zero with probability $1 - \delta$, that is,*

$$\mathbb{P}\left(\lim_{t \rightarrow \infty} \|\beta_t^s - \beta^{s*}\|^2 = 0\right) = 1 - \delta, \quad (2)$$

where the probability is over the random realizations of random samples $\{\Theta_t\}_{t=0}^{\infty}$. Furthermore, for the specific step size $\eta_t = \eta_0/(t + T_0)$ for some constants η_0 and T_0 , the model parameters at iteration t satisfy

$$\mathbb{E}_{\Theta} [f(\beta_t^s, \Theta)] - \mathbb{E}[f(\beta^{s*}, \Theta)] \leq \frac{C_0}{T_0 + t}, \quad (3)$$

with probability $1 - \delta$. Here, C_0 is a constant depending on the parameters of the sketching scheme, the above assumptions, and the objective function.

simulations

$$\mathbf{x}_i \sim \mathcal{N}(0, 1)$$

$$y_i = \mathbf{x}_i \boldsymbol{\beta}^*$$

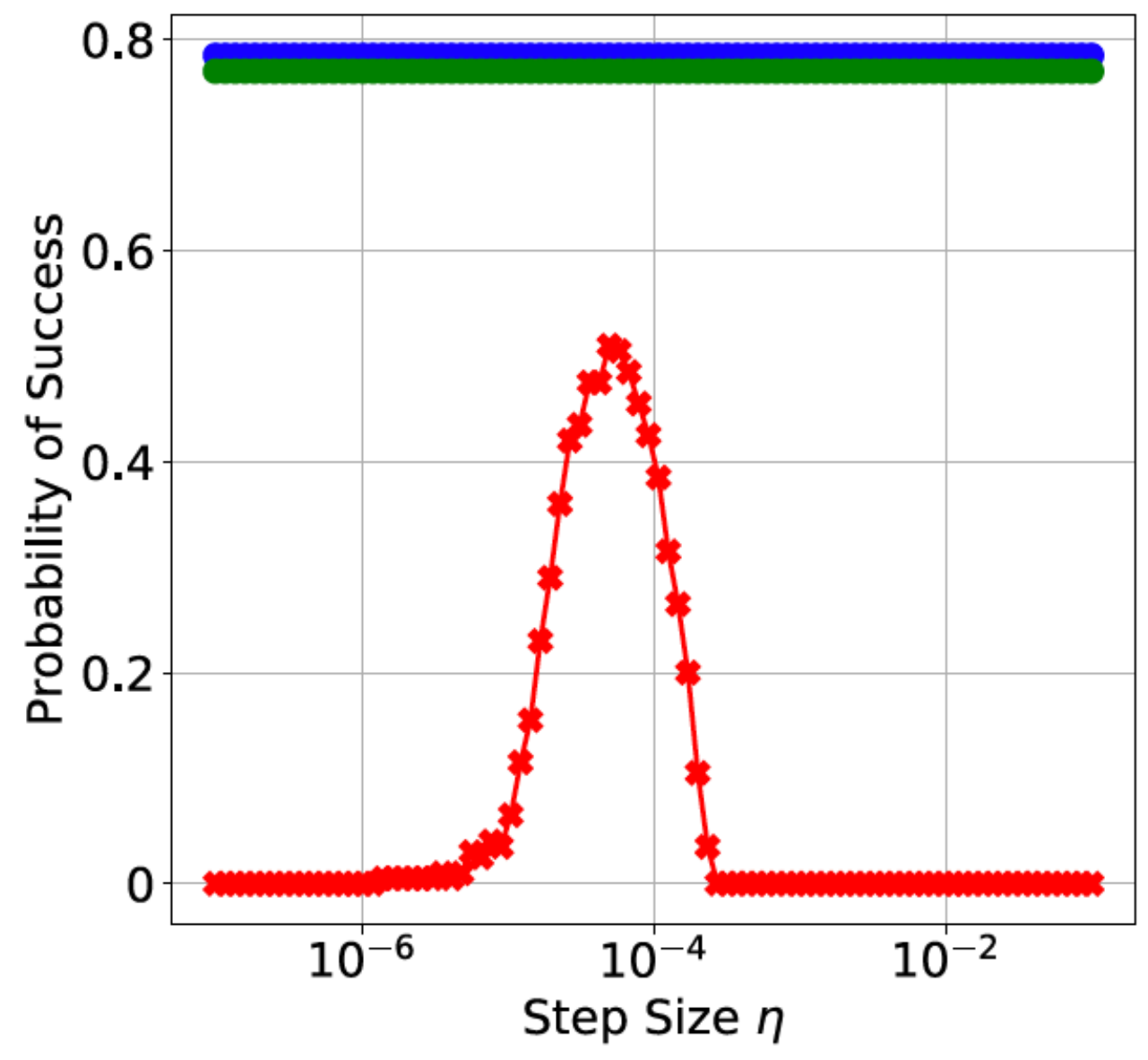
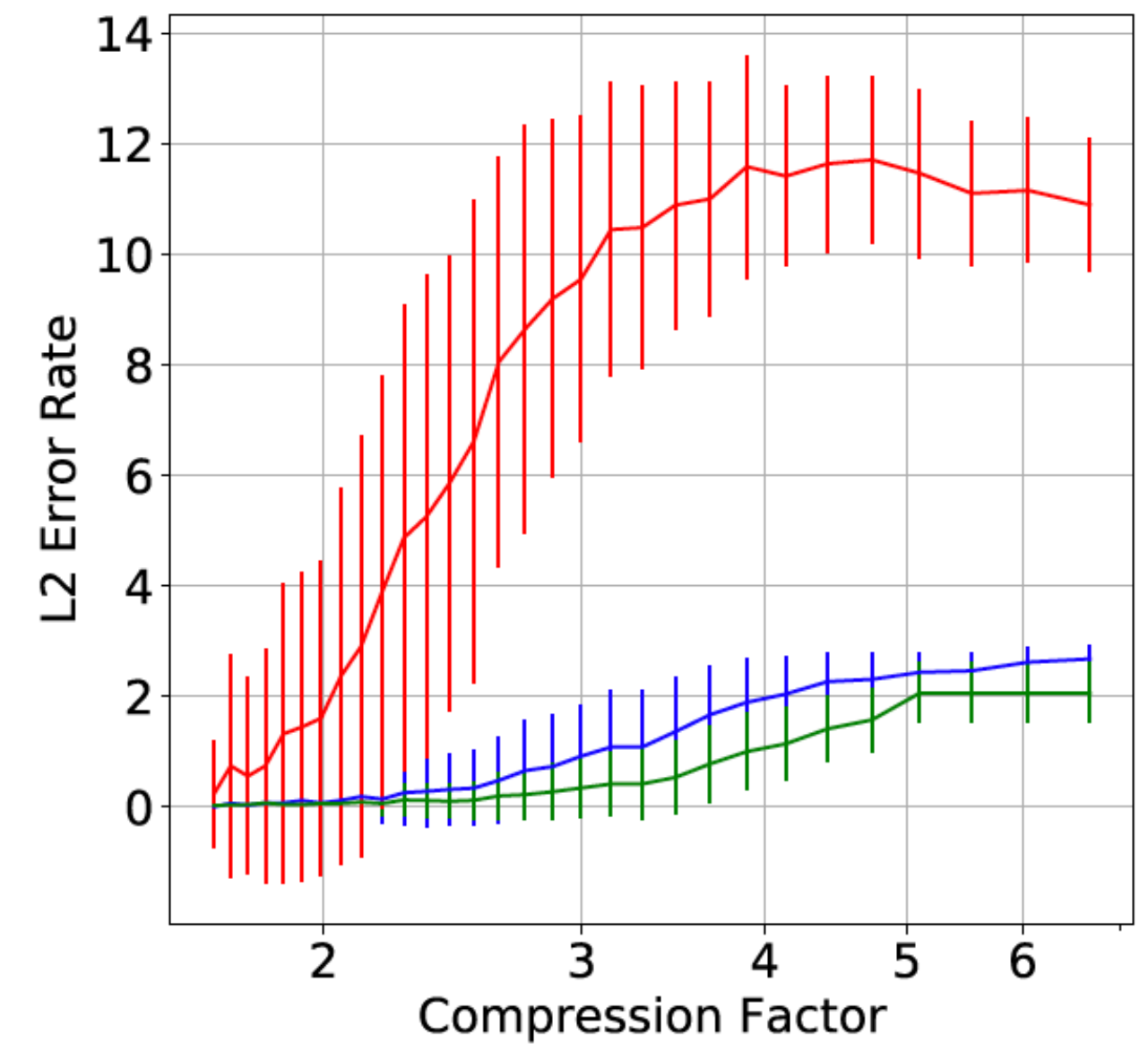
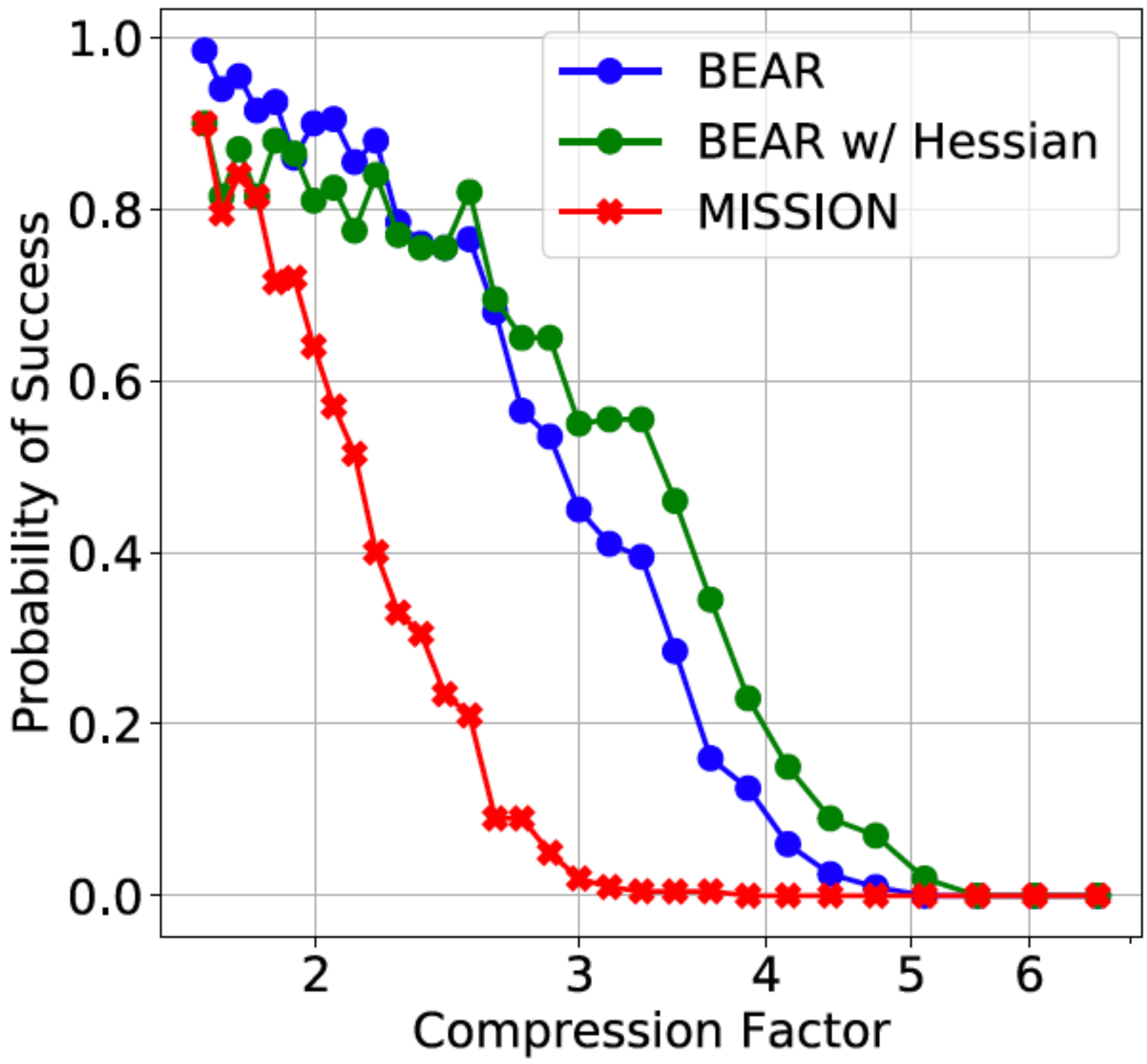
$\boldsymbol{\beta}^* : k - \text{sparse}$

$$\text{CF} = \frac{p}{\text{size of CS}}$$

$$p = 1000$$

$$n = 900$$

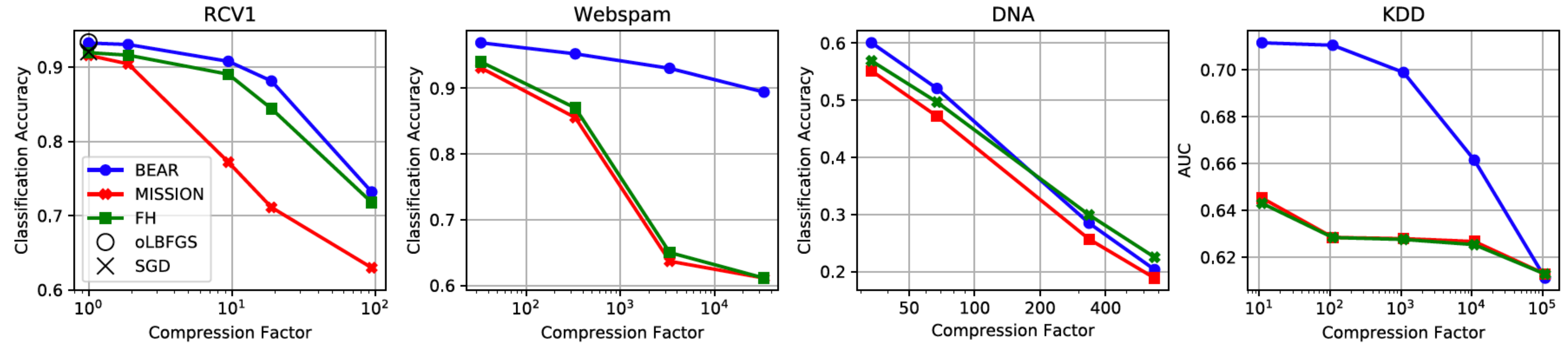
$$k = 8$$



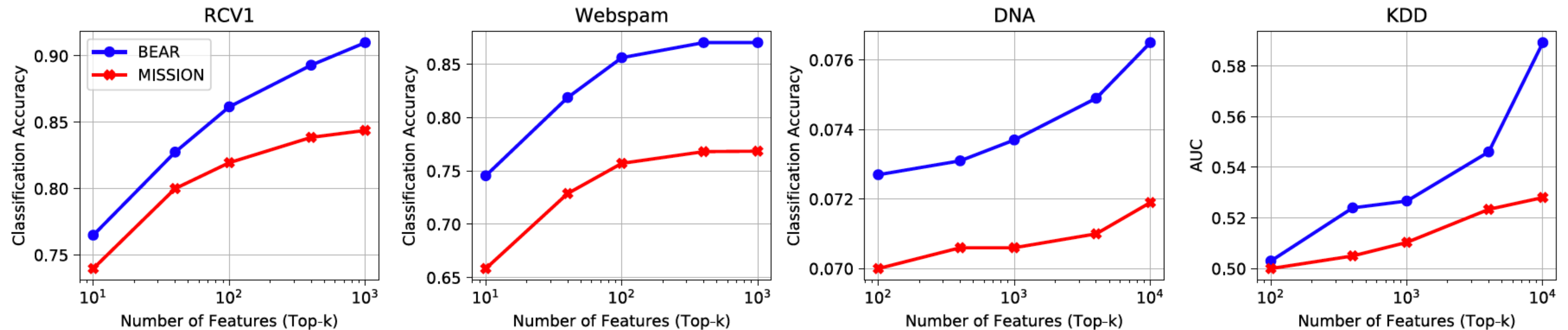
real-world experiments

Data set	Dim (p)	#Train (n)	#Test	Size	#Act.
RCV1	47,236	20,242	677,399	1.2GB	73
Webspam	16,609,143	280,000	70,000	25GB	3730
DNA	16,777,216	600,000	600,000	1.5GB	89
KDD 2012	54,686,452	119,705,032	29,934,073	22GB	12

Classification



Feature selection



summary and future directions

- **adaptively** learn the hashing scheme in the Count Sketch based on the stochastic gradients
- efficient training of massively large **nonlinear models** using LBFGS + sketching (Transforms, etc.)
- **distributed** learning/analysis using LBFGS + sketching explore communication-computation tradeoff

Thanks!

- find the **paper** at <https://arxiv.org/abs/2010.13829>
- find the **code** at <https://github.com/BEAR-algorithm/BEAR>